SSL/TLS transport protocol

SSL / TLS

APM@FEUP

The SSL/TLS protocol

> Web traffic (as all network traffic) is subject to many threats

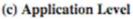
- integrity, confidentiality, authentication theft, denial of service, ...
- need added security mechanisms

Traffic security can appear at several levels in the network

- At the lowest protocol level (IPSec) embedded in the network
- Just above TCP level but used by several high-level protocols
 - Implemented in specific packages (e.g., browsers and web servers)
- At the application level using underlying standard protocols

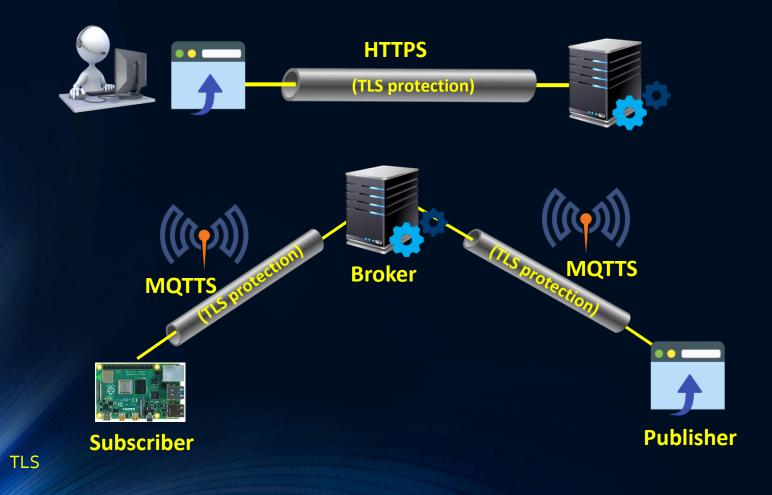
			нттр	FTP	SMTP	
HTTP	FTP	SMTP	5	SL or TL	s	Kerberos
	тср			TCP		
IP/IPSec		IP				
(a)	(a) Network Level			ransport	Level	(c)

	S/MIME			
Kerberos	SMTP	HTTP		
UDP	TCP			
IP				



TLS as the underlying protocol

TLS can be the security mechanism added to other higher-level protocols like HTTP or MQTT



SSL / TLS evolution and standards

SSL (secure sockets layer) is a transport protocol

- Originally developed by Netscape
- Makes use of TCP to provide a reliable end-to-end service
- Version 3 was presented as a draft internet standard

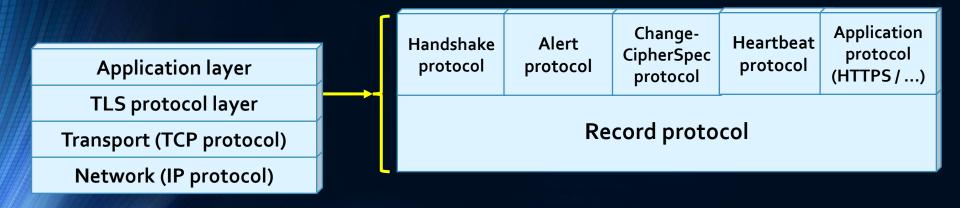
The approval of SSL by IETF became the TLS standard

- TLS Transport Layer Security (RFC 2246 TLS 1.0, 1999)
- TLS 1.0 is essentially SSL v. 3.1 and was backward compatible not causing disruptions
- TLS evolved with standards RFC 4346 in 2006 (1.1) and RFC 5246 in 2011
 - Improved cryptography (e.g., SHA-256 and AES, better encryption modes, ...)
- RFC 6167 (2011) and RFC 7568 (2015) refined all TLS versions
 - removed backward compatibility with SSL 2.0 and SSL 3.0
- TLS 1.3 (RFC 8446) was already approved (2018)
 - operations independent of cipher suites
 - removing support for weaker and lesser used cryptography algorithms
 - Session hash and new signature and key exchange algorithms

TLS stack and architecture

TLS protocol is composed of two layers of sub-protocols

- Handshake allows encryption, MAC and keys negotiation and authentication
- Alert conveys alert messages and errors
- ChangeCipherSpec allows updating the cipher suite in use and making it current
- Heartbeat checks link operation and prevents disconnection
- The Record protocol is the data transmission format for exchanging application data



TLS architecture and properties

The Handshake protocol establishes a TLS session

- Association and link between client (browser) and server
- Defines a set of cryptographic parameters
- Conducts a server authentication in the client using a certificate
- Optionally can do the same relatively to the client (seldom used)

The TLS secure channel provides three properties

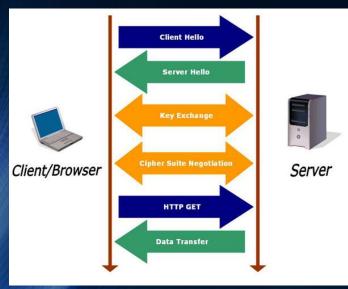
- Confidentiality using symmetric encryption and decryption with keys generated for the session
- Integrity of data using a message hash and a MAC also with keys generated for the session
- Server authentication preventing server spoofing and the establishment of connections to attackers unknowingly
 - Using a certificate for the server domain or organization

Handshake and data transfers

Initiated by the client to establish a TLS session

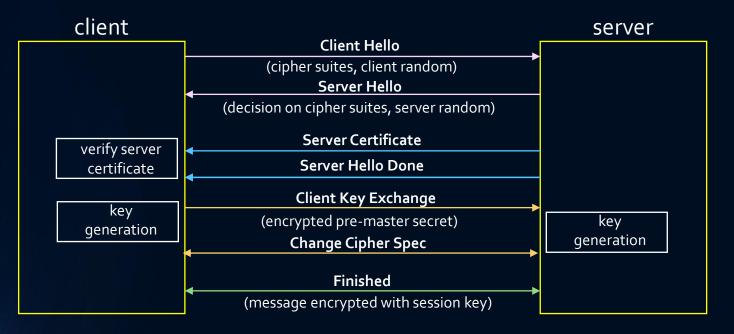
- It has several phases and consists of the exchange of simple messages
- negotiate encryption, mode, hash and MAC algorithms
- authenticate the server (optionally also the client)
- generate and exchange session cryptographic keys

The keys established in the handshake protocol are used for HTTP requests through the Record protocol



1 byte	3 bytes	≥ 0 bytes					
Туре	Length	Content					
Handshake Protocol							

Handshake operations



Client Hello – transmits the highest version understood and supported algorithms (crypto algorithms, key exchange, and compression) in preferred order. Also, a value (client_random) composed of a time stamp and a cryptographic random value is transmitted. Server Hello – transmits the decision of the server concerning version and algorithms. A similar server_random value is also transmitted.

The server must send his certificate, and its thorough verification in the client is crucial. The Server Hello Done message terminates the algorithms negotiation

The Client Key Exchange message transmits a secret (pre-master-key) generated in the client and usually encrypted using the public key of the server certificate. This key is used by both client and server, together with the client and server_random values, to generate the session keys.

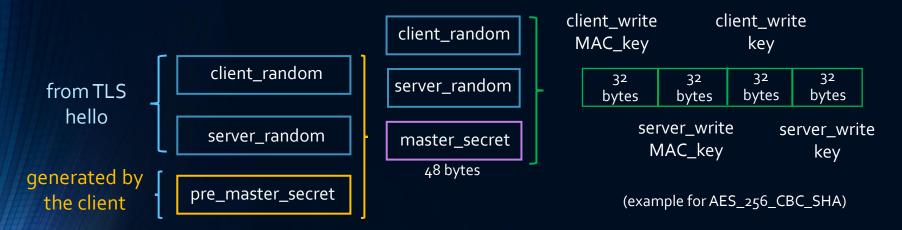
The Change Cipher Spec messages turn the session keys effective.

Finally, the Finished messages uses hash, MAC and encryption to test the session keys in both senses.

TLS session keys generation

From the pre-master-secret both sides gain the same keys

- The pre-master-key is generated by the client and transmitted to the server (using RSA or DH)
 - Then a master-secret is computed in both sides
 - Finally, the session keys are computed in both sides



The master_secret and session keys are generated in both sides using an agreed upon pseudo-random function

 $u_0 = label || server_random || client_random, where || denotes concatenation and label a string like "master secret"$ $<math>u_i = HMAC(secret, u_{i-1}), where secret is the premaster or master secret output = u_1 || u_2 || ..., retaining only the necessary bytes$

TLS Data Transmission

Data transmission is in both directions using records

Follows the TLS Record protocol data format

• Is used by all the TLS sub-protocols, after key exchange and generation phase in the Handshake protocol

Contains a header and a payload

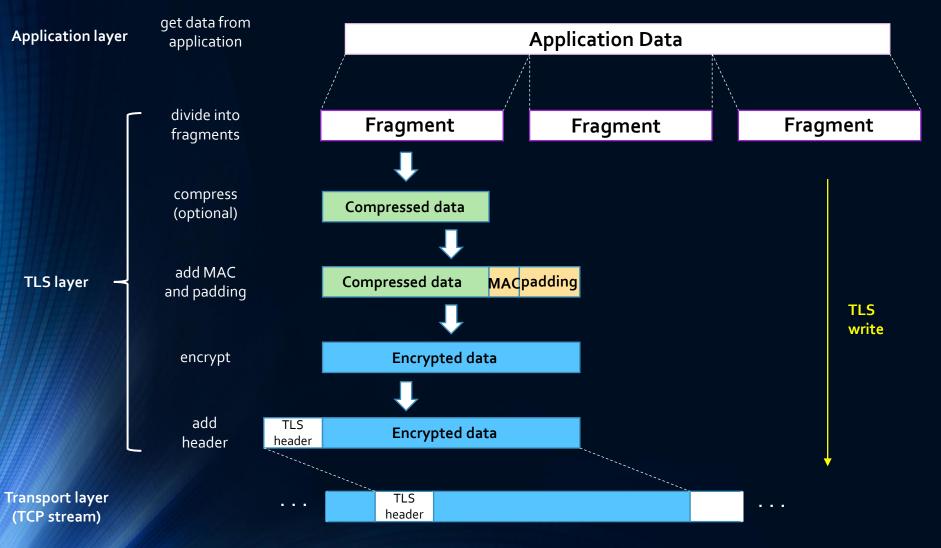
- The <u>header</u> is divided into
 - <u>Content type</u> (1 byte) indicating the sub-protocol (Handshake, ChangeCipherSpec, Alert, Heartbeat and Application
 - <u>Version</u> (2 bytes) indicating the SSL / TLS version
 - <u>Length</u> (2 bytes) with the payload length in bytes, until a maximum of 2¹⁴
- The <u>payload</u> contains the data transmitted in this record (may be compressed), with an appended MAC code and padding, all encrypted
 - The message presented to the MAC algorithm contains the record sequence number, the compression type and version (if any), the length (after possible compression), and the compressed bytes (fragment)

TLS Record

Content type	Version	Length	Data (may be compressed)		Padding
TLS header TLS		er	TLS payload (encrypted)		10

TLS write operation

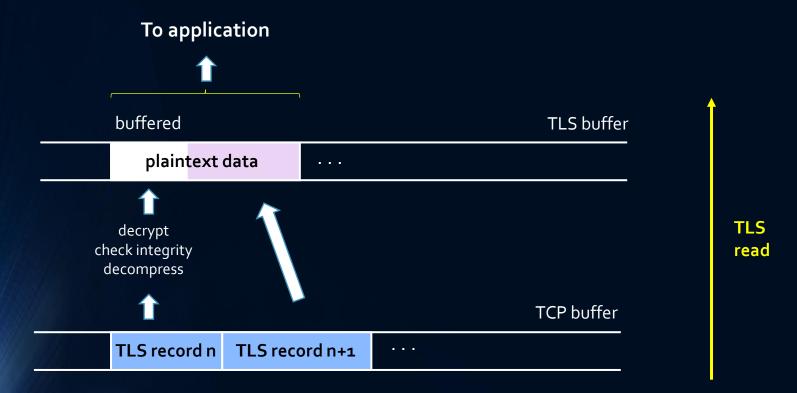
Usually, data from HTTP requests or responses (application)



TLS read operation

The plaintext data is recovered before sent to application

the browser or web server showing and running the web application



HTTPS

HTTP over TLS

- Combination of HTTP and TLS to secure communications between browser and server
- follows the IETF standard RFC 2818
 - specifies TLS handshake followed by normal HTTP requests and responses
 - no fundamental changes from SSL to TLS
- The URL begins with https://... rather then http://...
- Uses port 443 instead of port 80 (by default)
- Allows confidentiality and integrity over the HTTP data
 - URL addresses
 - document contents
 - form data
 - cookies
 - HTTP headers